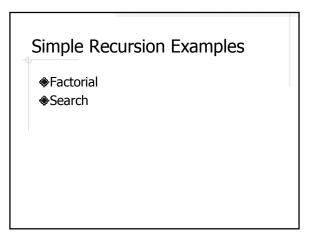


Ending Condition ♦ When the recursion should stop ♦ To avoid infinite recursion, make sure the ending condition n Exists n Reachable n Comes before the recursive call



String Permutation

- Output all the permutations of n characters
 - m E.g. "abc"
 wabc, acb
 wbac, bca
 wcab, cba
- How do we reduce the problem of n characters to the problem of n-1 characters??

Fibonacci Series

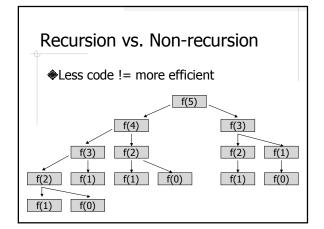
- **♦**0, 1, 1, 2, 3, 5, 8, 13, 21, ...
- Definition
 - $_{n}$ fibonacci(0) = 0
 - $_{n}$ fibonacci(1) = 1
 - n fibonacci(n) = fibonacci(n-1)+fibonacci(n-2)

Recursive Fibonacci

```
int fibonacci( int n )
{
  if( n == 0 ) return 0;
  else if( n==1 ) return 1;
  else
    return fibonacci(n-1) + fibonacci(n-2);
}
```

Non-recursive Fibonacci

```
int fibonacci( int n )
{
    if( n == 0 || n==1 ) return n;
    int last1 = 1, last2 = 0, fibo;
    for( int i=2; ??; ++i )
    {
        fibo = last1+last2;
        ??
    }
    return fibo;
}
```



Timing

- The best way to appreciate a good algorithm is to see how fast it runs
- And time it
- \$System.currentTimeMillis()
- ♦System.nanoTime()

When Can We Use Recursion?

- A problem itself is recursively defined
 - Fibonacci f(n) = f(n-1) + f(n-2)
 - n Tree
 - w A tree has a root
 - $\mbox{\sc w}$ Each child of the root is also a tree
- A problem of size n can be reduced to a problem of size less than n
 - ր Factorial: n n-1
 - n Sort: n n-1
 - _n Binary search: n n/2

When Should We Use Recursion?

- ♦ When the homework problem says so
- When speed of code development takes precedence over code efficiency
- When the problem is naturally recursive Fibonacci Series
- When the non-recursive solution is much harder
 - n Hanoi tower
 - n Solving maze